

## Haskell Live

# [07] Aufgabenblatt 4 (Bäume) und “What the $\$(.) \backslash x \rightarrow x?!$ ”

Bong Min Kim

e0327177@student.tuwien.ac.at

Christoph Spörk

christoph.spoerk@inode.at

Florian Hassanen

florian.hassanen@gmail.com

Bernhard Urban

lewurm@gmail.com

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## What the $\$(.) \backslash x \rightarrow x?!$

```
-- ($), (.) und flip sind bereits in Prelude
-- definiert, wir wollen sie aber selbst definieren
import Prelude hiding (($), (.), flip)
import Data.Char
import Data.List
foolist :: [Integer]
```

```
foolist = [1..1000]
uselessfkt0 :: Integer -> [Integer] -> Integer
uselessfkt0 x l = sum (filter tollespraedikat l)
  where tollespraedikat y = (y `mod` x) == 0
```

## Lambda Expressions

```
-- in hugs: > (\x -> x + x) 5 = 10
-- soll an  $\lambda x \rightarrow x + x$  erinnern
myadd :: Integer -> Integer -> Integer
myadd = (\x y -> x + y)
-- in hugs: myadd 2 4 = 6
uselessfkt1 :: Integer -> [Integer] -> Integer
uselessfkt1 x l = sum (filter (\y -> (y `mod` x) == 0) l)
-- oder gesamte Funktion als Lambda Expression
uselessfkt2 :: Integer -> [Integer] -> Integer
uselessfkt2 = \x l -> sum (filter (\y -> (y `mod` x) == 0) l)
```

( $\$$ )

```
-- aus der Prelude Definition: Applikationsoperator
( $\$$ ) :: (a -> b) -> a -> b
f $ x = f x
-- schwächste Priorität für ( $\$$ )
infixr 0 $
-- => Klammern sparen! (Lisp hat uns ohnehin schon zu viele gekostet)
uselessfkt3 :: Integer -> [Integer] -> Integer
```

```

uselessfkt3 x l = sum $ filter (\y → (y `mod` x) ≡ 0) l
-- besonders bei vielen Funktionsapplikationen praktisch...
yauf1, yauf2 :: String
yauf1 = take 4 (snd (splitAt 6 ("for" ++ "teh" ++ "lulz" ++ "haha")))
yauf2 = take 4 $ snd $ splitAt 6 $ "for" ++ "teh" ++ "lulz" ++ "haha"

```

## flip

```

-- aus der Prelude Definition
flip :: (a → b → c) → b → a → c
flip fkt x y = fkt y x
subtract :: Int → Int → Int
subtract x y = (flip (-)) x y

```

## (.)

```

-- aus der Prelude Definition: Funktionskomposition
(◦) :: (b → c) → (a → b) → (a → c)
(f ◦ g) x = f (g x)
-- stärkste Priorität für (.)
infixr 9 ◦
myToUpper1, myToUpper2, myToUpper3 :: Char → Char
myToUpper1 ch = chr $ ord ch - 0 x20
myToUpper2 ch = (chr ◦ (subtract 0 x20) ◦ ord) ch
uselessfkt4 :: Integer → [Integer] → Integer
uselessfkt4 x l = (sum ◦ filter (\y → (y `mod` x) ≡ 0)) l

```

## Pointfree

```
myToUpper3 = chr ∘ (subtract 0 x20) ∘ ord
-- on the way to pointfree...
uselessfkt5 :: Integer → [Integer] → Integer
uselessfkt5 x = sum ∘ filter (λy → (y `mod` x) ≡ 0)
-- pointfree (thx @lambdabot ;-))
uselessfkt6 :: Integer → [Integer] → Integer
uselessfkt6 = (sum ∘) ∘ filter ∘ flip flip 0 ∘ ((≡) ∘) ∘ flip mod
```

Expression	Typ
(sum .)	(b -> [a]) -> b -> a
filter	(a -> Bool) -> [a] -> [a]
flip	(a -> b -> c) -> (b -> a -> c)
flip flip	a -> (b -> a -> c) -> (b -> c)
flip flip 0	(b -> a -> c) -> (b -> c)
(.)	(a -> b) -> (c -> a) -> c -> b
(==)	a -> a -> Bool
((==) .)	(b -> a) -> b -> a -> Bool
flip mod	a -> a -> a

⇒ Pointfree ist nicht immer sinnvoll! Für Interessierte: <http://www.haskell.org/haskellwiki/Pointfree>

Hausaufgabe: Wo könnte der Pointfreestyle bei Aufgabe6 sinnvoll sein?

## Aufgabenblatt 4

```
-- my tree definition
data Tree = Leaf Integer |
  Node Integer Tree Tree deriving Show
```

```

type Layer = [Integer]
data MyOrd = BottomUp | TopDown

  -- some trees
t1 = (Node 5 (Node 5 (Leaf 4) (Leaf 2)) (Leaf 3))
t2 = (Node 5 (Node 5 (Leaf 4) (Leaf 2)) (Node 3 (Leaf 1) (Leaf 3)))
t3 = (Node 5 (Node 5 (Leaf 4) (Node 2 (Leaf 1) (Leaf 3))) (Node 5 (Node 5 (Leaf 4) (Node 2 (Leaf 1) (Leaf 1))) (Leaf 3)))
t4 = (Node 1 (Node 1 (Node 1 (Leaf 2) (Leaf 3)) (Node 1 (Leaf 2) (Leaf 3))) (Node 1 (Node 1 (Leaf 2) (Leaf 3)) (Node 1 (Leaf 2) (Leaf 3)))

  -- a
mergeLayer :: [Layer] → [Layer] → [Layer]
mergeLayer [] r = r
mergeLayer l [] = l
mergeLayer (x1 : x1s) (x2 : x2s) = (x1 ++ x2) : (mergeLayer x1s x2s)
writeLayer :: Tree → MyOrd → [Layer]
writeLayer (Leaf x) _ = [[x]]
writeLayer (Node x t1 t2) TopDown = [x] : merged
  where merged = mergeLayer (writeLayer t1 TopDown) (writeLayer t2 TopDown)
writeLayer t BottomUp = reverse $ writeLayer t TopDown

  -- b
data STree = Nil |
  SNode Integer STree STree deriving Show
treeToSortedList :: Tree → [Integer]
treeToSortedList t = sort $ nub $ foldr (++) [] (writeLayer t TopDown)
splitHalf :: [a] → ([a], a, [a])
splitHalf l = ((take p l), (l !! (p)), (drop (p + 1) l))
  where p = (length l) `div` 2
listToStree :: [Integer] → STree
listToStree [] = Nil
listToStree l = SNode x (listToStree l1) (listToStree l2)
  where (l1, x, l2) = splitHalf l

```

```

transform :: Tree → STree
transform t = listToStree (treeToSortedList t)
    -- some tree functions
    -- calculate tree depth
depth :: Tree → Integer
depth (Leaf _) = 0
depth (Node _ subt1 subt2) = 1 + (max (depth subt1) (depth subt2))
flatten :: Tree → [Integer]
flatten (Leaf x) = [x]
flatten (Node x subt1 subt2) = (x : ((flatten subt1) ++ (flatten subt2)))
treemap :: (Integer → Integer) → Tree → Tree
treemap f (Leaf x) = Leaf (f x)
treemap f (Node x subt1 subt2) = Node (f x) (treemap f subt1) (treemap f subt2)
    -- tree printer
space x = map (\x → ' ' ) [1..x]
zeroCopy :: Tree → Tree
zeroCopy (Leaf _) = (Leaf 0)
zeroCopy (Node _ subt1 subt2) = (Node 0 (zeroCopy subt1) (zeroCopy subt2))
setRoot :: Integer → Tree → Tree
setRoot r (Node _ subt1 subt2) = (Node r subt1 subt2)
    -- transformers
    -- balance fills up the given binary tree to full binary tree
balance :: Tree → Tree
balance (Leaf x) = Leaf x
balance (Node x s1@(Node y subt1 subt2) (Leaf z)) = (Node x b1 (setRoot z (zeroCopy b1)))
    where b1 = balance s1
balance (Node x (Leaf z) s2@(Node y subt1 subt2)) = (Node x (setRoot z (zeroCopy b2)) b2)
    where b2 = balance s2
balance (Node x subt1 subt2)

```

```

| bal1d > bal2d = (Node x bal1 (balance_ bal2 bal1))
| bal1d < bal2d = (Node x (balance_ bal1 bal2) bal2)
| otherwise = (Node x bal1 bal2)
where bal1 = balance subt1
        bal2 = balance subt2
        bal1d = depth bal1
        bal2d = depth bal2

-- the first tree gets the same structure as the second one
balance_ :: Tree → Tree → Tree
balance_ (Leaf x1) (Leaf x2) = (Leaf x1)
balance_ (Node x s1 s2) (Leaf _) = (Node x s1 s2)
balance_ (Leaf x) b1@(Node y s1 s2) = (setRoot x $ zeroCopy b1)
balance_ (Node x1 s1 s2) (Node _ s3 s4) = (Node x1 (balance_ s1 s3) (balance_ s2 s4))

mergeTreeShow [] [] = []
mergeTreeShow (t1 : t1s) (t2 : t2s)
  | (l `mod` 2) ≡ 0 = ((t1 ++ (space 5) ++ t2) : (mergeTreeShow t1s t2s))
  | otherwise = ((t1 ++ (space 4) ++ t2) : (mergeTreeShow t1s t2s))
where
  str = (t1 ++ (space 5) ++ t2)
  indices = findIndices (≠ ' ') str
  index1 = indices !! 0
  index2 = indices !! 1
  l = index2 - index1

treeshow_ :: Tree → [String]
treeshow_ (Leaf x) = [show x]
treeshow_ (Node x subt1 subt2) = (help (head mt) x) ++ mt
  where t1 = treeshow_ subt1
        t2 = treeshow_ subt2
        mt = mergeTreeShow t1 t2
        help s x
          | ((index2 - index1) > 2) ∧ (ch ≡ '0' ∨ ch ≡ '- ') = (help str2 x) ++ [str2]

```

```

| ((index2 - index1) > 2) ∧ x ≠ 0 = (help str1 x) ++ [str1]
| otherwise = [(space (index1 + 1)) ++ (show x) ++ (space (l - index2))]
where
  indices = findIndices (≠ ' ') s
  index1 = indices !! 0
  index2 = indices !! 1
  ch = s !! index1
  l = length s
  str1 = (space (index1 + 1)) ++ [ '/' ] ++ (space (index2 - index1 - 3)) ++ [ '\\ ' ] ++ (space (l - (index2)))
  str2 = (space (index1 + 1)) ++ [ '- ' ] ++ (space (index2 - index1 - 3)) ++ [ '- ' ] ++ (space (l - (index2)))
treeShow :: Tree → IO ()
treeShow t = sequence_ (map putStrLn (map (map pp) (treeShow_ (balance t))))
  where
    pp c
      | c ≡ '0' ∨ c ≡ '- ' = ' '
      | otherwise = c
treeOrigshow :: Tree → IO ()
treeOrigshow t = sequence_ (map putStrLn (treeShow_ $ balance t))

```